

For this TP you can start from the file `tp3_code.py` which contains the LFSR functions and Berlekamp-Massey.

## Cryptanalysis of the Geffe Cipher

The Geffe cipher is a stream cipher with 3 combined registers, proposed by Geffe in 1973. It combines three binary LFSRs using the Boolean function:

$$f(x_1, x_2, x_3) = x_3 + x_2x_3 + x_1x_2 \quad .$$

In this exercise, we suppose that the internal LFSRs are the following:

- LFSR1 of length 13, polynomial  $P_1 = 1 + X + X^3 + X^4 + X^{13}$
- LFSR2 of length 11, polynomial  $P_2 = 1 + X^2 + X^{11}$
- LFSR3 of length 9, polynomial  $P_3 = 1 + X^4 + X^9$

During initialization, the three LFSRs are initialized with their respective initial states  $S_1, S_2, S_3$ . We note respectively  $s_1(t), s_2(t), s_3(t)$  the output bits at time  $t$ , and  $z(t)$  the output bit of the combined LFSR.

**Question 1.** *What is the value of  $z(t)$  depending on  $s_1(t), s_2(t)$  and  $s_3(t)$ ?*

**Question 2.** *Program in Python a function `geffe(S1,S2,S3,N)` that takes as input the three initial states, and an integer  $N$ , and returns the  $N$  first bits of the sequence.*

Check that the 20 first bits of the sequence generated with the generator, initialized with the states:

$$\begin{cases} S_1 = [1, 0, 1, 0, 1, 0, 1, 0, 1, 0, 1, 0, 1] \\ S_2 = [1, 0, 1, 0, 1, 0, 1, 0, 1, 0, 1] \\ S_3 = [1, 0, 1, 0, 1, 0, 1, 0, 1] \end{cases}$$

are:

$$Z = 1, 0, 1, 0, 1, 0, 1, 0, 1, 1, 1, 0, 1, 0, 0, 0, 0, 1, 0, 1, \dots$$

**Question 3.** *What is the linear complexity of each internal LFSR? In theory, what is the linear complexity of the Geffe generator? Check with the Berlekamp-Massey algorithm.*

## Correlation Attack

**Question 4.** *What is the complexity of exhaustive search to determine the initial state of the Geffe generator?*

The goal of this section is to describe the *correlation attack* proposed by Siegenthaler in 1985 on the Geffe generator.

We assume that at each time  $t$ , the bits  $s_1(t), s_2(t), s_3(t)$  produced by the internal LFSRs are independent random variables, uniformly distributed in  $\mathbb{F}_2$ . Then,  $z(t) = f(s_1(t), s_2(t), s_3(t))$  is also a random variable taking values in  $\mathbb{F}_2$ .

**Question 5.** *Show that:*

$$\Pr[z(t) = s_1(t)] = \Pr[z(t) = s_3(t)] = \frac{3}{4}$$

and

$$\Pr[z(t) = s_2(t)] = \frac{1}{2}$$

**Question 6.** Let  $x$  be a random variable uniformly distributed over  $\mathbb{F}_2$  and independent from  $z(t)$ . What is the probability:  $p = \Pr[z(t) = x]$ ?

We have shown that the output of the combined LFSR is strongly correlated with the output of the two LFSRs LFSR1 and LFSR3. The goal of a correlation attack is to exploit this property to deduce the initial states of each register.

We suppose that using a known plaintext, we have determined  $\ell$  bits of the sequence, from  $z(t_0)$  to  $z(t_0 + \ell - 1)$ .

**Finding the internal state of LFSR1.** We do an exhaustive search on the internal state of LFSR1 at time  $t_0$ , and for each candidate state  $\tilde{S}_1$ , we compute the sequence  $s_1(t)$ . When the internal state is correct, we expect that  $(s_1(t))$  coincides with  $\simeq 3/4$  of the values if  $\ell$  is large enough. Otherwise, we expect that only  $\ell/2$  elements coincide.

**Finding the internal state of LFSR3.** We do the same with LFSR3, and we obtain a candidate  $\tilde{S}_3$  for its internal state at time  $t_0$ .

**Finding the internal state of LFSR2.** Finally we perform an exhaustive search for the internal state of LFSR2. If the state is correct we expect to obtain exactly the output sequence.

**Question 7.** What is the complexity of the attack to find  $S_1$  and  $S_3$ ? Of the whole attack? Is it better than exhaustive search?

**Question 8.** Implement this attack. The file `tp3_code.py` contains a sequence of 100 bits generated by the Geffe generator (also given below). Find the corresponding initial state of the three LFSRs.

Challenge sequence:

```
challenge = [0, 0, 0, 0, 0, 0, 0, 1, 0, 1, 0, 0, 1, 0, 0, 0, 1, 0, 0, 0, 1,
0, 0, 0, 0, 0, 1, 0, 0, 0, 0, 0, 1, 1, 1, 1, 1, 0, 1, 1, 1, 1, 0, 0,
0, 0, 1, 0, 0, 1, 1, 1, 1, 0, 0, 0, 0, 0, 1, 0, 0, 1, 1, 0, 1, 1, 0,
1, 1, 1, 0, 1, 0, 1, 1, 0, 1, 0, 1, 0, 1, 0, 0, 1, 1, 1, 0, 0, 0, 0,
0, 0, 0, 0, 0, 1, 0, 0, 1, 0]
```