Security Properties of Hash Functions

Let $h : \{0,1\}^* \to \{0,1\}^n$ be a hash function that we suppose collision resistant. Let h' be the following function:

$$h' : \begin{cases} \{0,1\}^* \to \{0,1\}^{n+1} \\ x \mapsto \begin{cases} 0 \|x \text{ if } |x| = n \\ 1 \|h(x) \text{ otherwise} \end{cases}$$

Question 1. Show that h' is collision resistant.

Question 2. Show that h' is not preimage resistant.

Collisions

During this exercise, I suggest to generate random integers using the function randrange from the package random: randrange(N) generates a uniform random integer from $\{0, 1, \ldots, N-1\}$. You may start from the file tp1_code.py.

We use the library hashlib from Python's standard library, which implements many hash functions (MD5, SHA-1, SHA-2, SHA-3). These constructions take as inputs objects of type bytes and not str. In particular, you need to use the .encode() function or directly construct a byte string with the prefix "b". Read the documentation¹ for more information. Here is an example:

```
import hashlib
import hashlib
sha2 = hashlib.sha256
print(sha2(b"Hello world").hexdigest())
print(sha2("Hello world".encode()).hexdigest())
```

Collision on Truncated Hash Function

Question 3. Choose a prefix (for example your first name) and implement the generic collision search algorithm using the birthday paradox, to find two strings s_1 and s_2 both starting with this prefix, so that the 32 first bits (i.e., 8 first characters in the hex string) of $SHA2(s_1)$ and $SHA2(s_2)$ are equal. For example:

```
Collision found !

Input 1: maxime15857573905157511205

Input 2: maxime13871373172309900626

sha2 ( input1 ) =

a4de129026e4f1b46270dc73772a14c26d90c3df19d2a040d347cc154d38c4f8

sha2 ( input2 ) =

a4de1290d4324581554e4804b53f01f95211371a4241386372502d571fc1e06c

SHA256 prefix ( first 32 bits ) : a4de1290
```

Question 4. Estimate roughly the time and memory complexities of your algorithm. Does it depend on the length of your first name?

¹https://docs.python.org/3/library/hashlib.html

Collisions with Small Memory

In the following, we identify an *n*-bit truncated hexadecimal hash to an integer between 0 and $2^n - 1$, and use the function sha2Trunc defined in tp1_code.py.

Let $H : \{0, 1, ..., 2^n - 1\} \rightarrow \{0, 1, ..., 2^n - 1\}$ be a function (for example the sha2Trunc function). Starting from a message X_0 , we define the sequence $X_{i+1} := H(X_i)$. Since it takes values in a finite set, it's necessarily periodic after some point.

We denote by c the length of the pre-period X_0, \ldots, X_{c-1} (the length of the tail of the ρ) and ℓ the length of the cycle, so that $X_0, \ldots, X_{c+\ell-1}$ are all distinct.

Floyd's cycle-finding algorithm is given in Algorithm 1. It defines another sequence $Y_i = X_{2i}$, i.e., $Y_0 = X_0$ and $Y_{i+1} = H(H(Y_i))$, and outputs an element x that belongs to the cycle.

Algorithm 1 Floyd's cycle-finding algorithm.
$1: x \leftarrow H(X_0)$
2: $y \leftarrow H(H(X_0))$
3: while $x \neq y$ do
4: $x \leftarrow H(x)$
5: $y \leftarrow H(H(y))$
6: end while

Question 5. Implement Floyd's algorithm with a uniformly random value for X_0 .

Question 6. Write an algorithm to find the length ℓ of the cycle.

If c > 0 and $\ell > 1$ one has:

$$H(X_{c-1}) = X_c = X_{c+\ell} = H(X_{c+\ell-1})$$

and by definition $X_{c-1} \neq X_{c+\ell-1}$. We will use this property to find a collision.

Question 7. Write two new chains that start from X_0 and X_{ℓ} and compute X_i and $X_{i+\ell}$ until equality. During the computation of the chains, you should remember the previous element (thus 4 variables are required), and output it when $X_i = X_{i+\ell}$. Check that you indeed obtained a collision.

Question 8. Deduce examples of collisions on the n first bits of SHA256 for n = 16, 32, 40, 48.